

Agnostic approach to resolve the synchronization problem indistributed heterogeneous WSN's

Gurpreet Singh¹

Ramgarhia Institute of Engineering and Technology, Phagwara

ABSTRACT : The Proliferation in Micro-Electro-Mechanical system technology, has lure Many researcher to contribute continuously and successfully implanted these many tiny hardware to solve and mitigate many hassle of human life in real world, variety of phenomena are still under contemplate to fulfill the real life requirement to motivate the human lifestyle. In future scope many middle ware approaches are also advocate to condemn the human interaction and automate the system independently. Adverse implementation and poor protocol design can lead to many resource damages and serious harm to the network performance, so required to balance all unrelenting approaches in software and hardware design of WSNs. Our search in this paper is to maintain the hardware clock tick so as to synchronized clock perfectly without any delay or ahead, indeed many conventional approaches are in used but all of are more inclined to resource utilization in turns lead to wastage of many energy source which can harm the network subsequently and not beneficial for prolong network life in heterogeneous environment.

KEYWORDS: WSN, Synchronization, precision,

I. INTRODUCTION

Extensive research have done to improve the measure of performance in WSN's network, one of the prominent divergence in WSN is time clock used in hardware, lack in hardware resonator and their resonance frequency there is a pertinent requirement of precision and Synchronized, between the nodes, protocols of WSN's are designed by many researcher but still hold many disadvantage. The dimension of agility required to synchronized a clock is the delay in the time due to the propagation between the nodes in sensor network, estimator parameter such as clock drift and clock speed tumble in many nodes subsequently unreliable for the network. In this paper we will discuss the various time synchronization protocols used for node communication by incorporating the approach of SFTP and RBS routing protocols together, various scheme had already emerge from past decade and mitigate many anomalies in the network, but energy conserve and QoS are still the biggest challenge for the WSNs, to condense this problem we have design the protocol to conserve the resource with minimum usages in turn result into holistic improvement in the wireless network. Many application of wireless sensor network required time synchronization protocols, application such as habitat monitoring, environment, object target tracking to estimate the trajectory of the object, time synchronization protocol plays a vital role to get the actual aggregated information from the sensor nodes and to ameliorate the conventional method, inculcate with noble approach of synchronization might be beneficial for many applications that required real time re-synchronization communication between the communication network only. Many other methods for time synchronization is either global or local time synchronization. Service may require many hardware to update its clock but not cost effective seems critical to available to the entire node for many application all the time [1]. In this modeling we search variant type of time synchronization protocol and studied various ameliorate in traditional protocols. Efficiently and effectively variant protocols are implementing in fixed node environment with similar parameters to distinguish the working and provide the comprehensive analysis of the research work. To prolong network life time spam we implement the design in heterogeneous environment that is few of the nodes with more power to accommodate the network with maximum duration [2]. The node cock is measured as the function of the hardware oscillator denoted by equation below, clock and give three basic types of information Time of day, time interval and frequency, May inconsistent due to environment changes.

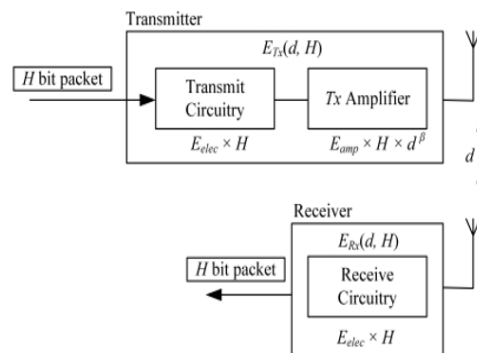
$$C(t) = K \int_{t_0}^t \omega(\tau) d\tau + C(t_0)$$

1

Where the $\omega(\tau)$ is the angular frequency and k is the constant, t is the time.

More specifically we target to coarse synchronization to many application in sensor network that required low precision synchronization in WSNs. When the network is deployed initial clock may not be running at same time and the quartz crystal at each node might be moving at different frequency gradually diverge from each other node, to inculcate the problem protocol level amelioration is indented to reduce the clock drift. Non deterministic transmission time cause by MAC layer also result into the hundreds of millisecond delay in the network, one approach is to broadcast the beacon signal to all the receiver that don't contain timing packets despite receiver will juxtapose their clock to compute relative offsets[3]

Figure 1.1 Energy Model of WSN



Every network in wireless network required a protocol to collect correct and accurate information among them one on us is Time Synchronization protocol; to manage the similar time between the nodes can be achieve by consistently maintaining the time. But many aspects of wireless network lead to decrease in the efficiency and accuracy of the network and become a challenge for us. Node intended to communicate in the field may generate a timestamp to transmit it to next neighbour node for synchronization, the packet hold the timestamp will introduce a variable amount of latency until it reaches to destination and is decoded at its respective receiver. The receiver may face the delay implicitly as compared the local clocks of the two nodes and found to topple in the accuracy. We can enumerate them into four major components as shown in figure 1.1, Send Time: The time needed to design a packet at the sender. The Overhead of operating system (such as context switches), and the time to transfer the message to the network interface for transmission. Access Time: Every packet experience delay at the Medium Access Control layer 2 in data link layer of TCP/IP model before actual transmission to avoid collision inside the network using expositional back off algorithm. The cause of such delay depend on the methods scheme used for channel access. Propagation Time: This is the total time required in propagation of the message between the network interfaces of the sender and the receiver. Receive Time: This is the time needed for the network of the receiver to receive the message and manage it to the host[3].

Even in fixed positions, nodes' communication ranges thus, topologies can change dramatically due to the vagaries of RF propagation, a result of its strong environmental dependence. These changes are difficult to predict in advance. Traditional large-scale networks such as the Internet work in the face of changing configurations and brittle software partly because the number of people maintaining the network has grown along with the size of the network itself[4]. In contrast, there may be a single human responsible for thousands of nodes in a dense sensor network. Any design in which each device requires individual attention is infeasible. This leads to another important requirement: sensor networks must be self-configuring, and adaptive to changes in their environment.

Yildirim and Kantarci[5], in their research has realized that the smaller the difference between the speeds of the clocks, the smaller the undesired effect of waiting times on the synchronization accuracy. They showed that the synchronization accuracy and scalability of slow-flooding can drastically be improved by employing a clock speed agreement algorithm among the sensor nodes. They presented an evaluation of this strategy on a test bed setup including 20 MICAz sensor nodes. They revealed that the smaller the error of the drift estimation and, hence, the difference between the speeds of the clocks, the smaller the undesired effect of waiting times on the synchronization accuracy. It was shown that the synchronization quality of slow-flooding-based time synchronization can drastically be improved by employing a clock speed agreement algorithm among the sensor nodes.

Motivation is one of the main requirements for research and it depends on need. In this section we focus on the factors and need that motivate us to work in the area. Clock synchronization is important for coordinating actions across a distributed set of actuators. Fuqiang Wang [6][7], proposed noble reliable time synchronization protocol to compensate the clock drift. In his research, authors try to opt for level fashion hierarchy communication, which holds same hop in the same level. Using linear regression and MAC time stamp try to condense the clock drift.

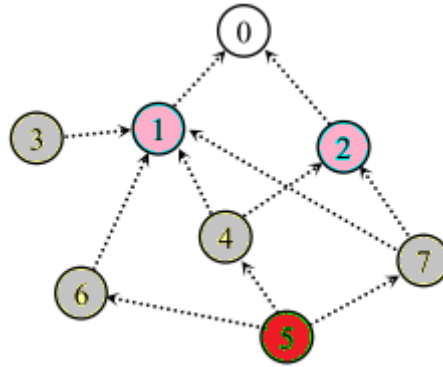


Figure1.3 random time source algorithms

Example of choice of time is exhibit on above figure1.3 whereas choice of time depends on the hops from root. Apparently node 1 and node2 can listen to the root 0 they are restricted to get synchronized with upper root node only similarly node with in the radio range of node 1 and node2 can make these two node potential time source. In same context node 5 can reliable on three potential time source node4, node6, node7 as portray in diagram above.

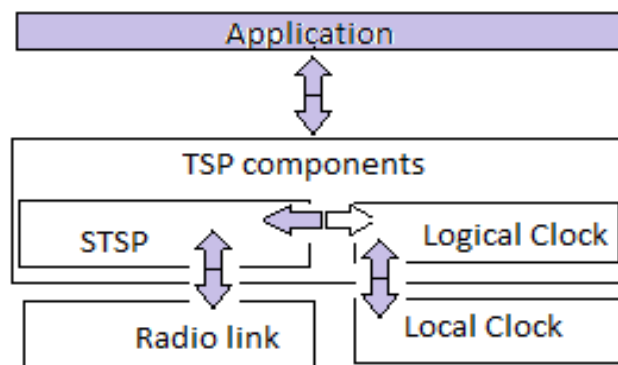


Figure 1.4 stack architecture of time synchronization service

In the above figure1.4 we exhibit Architecture of time synchronization service and its integration within the hardware and software platform.

Random selection minimized the chances of link failure and to re- synchronization of nodes as it can be easily comprehend that if any one of the parent hierarchy failure between node6, node4, node7 to child node5 still remain synchronized with other two, possibly consistent and robust in case of link failure[6]. In common parlance RTSP is a verified and tested protocol in ISA2420 platform, results reveals that the maximum average is 5 μ s and robustness in case of link failure.

II. LITERATURE REVIEWS

Jie Wu, Lyi Zhang, Yu Bai, and Yunshan Sun[1] proposed noble technique for time synchronization protocol using cluster consensus time synchronization on the basis of Distributed CTS algorithms so as to obtain fast synchronization between nodes for better energy performance. Protocol includes primary two part i.e inter and intra cluster synchronization. To get updated clock compensation parameter of network virtual and intracluster virtual clocks of every cluster head are assigned with corresponding weight based on the size of every cluster, the test show that proposed protocol in this literature have competency to improve the convergence rate due to combine technologies. Vikram Dhiman, Virpal Kaur, Abhijeet Singh [2] investigate the energy aware routing protocol for heterogeneous WSN, pragmatic routing scheme which hold hybrid and attributes for prolong the network life time. Proposed method involved many aspects and claiming multilevel chain hierarchy routing protocol manage to handle packet dissemination between source and destination. Later author investigate that the protocol design was efficient and adequate suitable for long distance Yosemite to provide cogent surveillance in natural disaster.

F. Sivrikaya and B. Yener [3] review some intrinsic properties of protocols such as energy, storage, computation and bandwidth. In this review problem of synchronization was explain explicitly and need for sensor network was expounded. Later researcher pin point the lack of infrastructure and used of sensor node as a medium to disseminate the message to destination and some motivational study. Elson, L. Girod, D. Estrin[4],

proposed a Synchronization scheme RBS novel technique to synchronous with third party instead of synchronization with sender and receiver, a reference beacon is used as time of arrival by receiving node to match the clock, author claim single broadcast received by all node at same time and propagation time that is equivalent to speed of light is relevantly very low can be neglected. RBS allow synchronization between m node using single pulse and number of beacon pulse to achieve precision. Kasim Sinan[5], proposed a Time Synchronous based on Slow flooding in wireless sensor network. The proposed method introduces accuracy and efficient time synchronous as compare to FTSP (Flooding Time Synchronization protocol). In traditional protocol common strategy is to flood current time information of a reference node into the network. Alternate method of same is proposed which allow node to propagate as soon as possible but this scheme is difficult have many limitation such as energy and network performance, in this paper author proposed a method to reduce the undesired effect of on synchronization accuracy without altering the propagation speed of the flood, main objective is to enhance the accuracy and stability of protocol SFTP drastically.

Fuqiang Wang, Peng Zeng, Haibin Yu and Xiaoquan Zhao[6][7], argues for the important aspects reliability of WSN's. Author proposed a new protocol to inculcate the Reliability named reliable Time Synchronization Protocol (RTSP) which is designed to adapt topology changes due to link failures or node mobility. Protocol was designed to work in level manner, which means nodes in each level must have same nodes. Protocol later designed and implemented for SIA2420 platform and use of TinyOS. Extensive literature work has been carried out and found that proposed scheme is more batter than traditional methods, results exhibits are gradually prominent and higher in the scale as compare to the traditional methods. Compensation mechanism and random time source opt to make the synchronization reliable and robust in case of node failure. Kasim Sinan[5], proposed a Time Synchronous based on Slow flooding in wireless sensor network. The proposed method introduces accuracy and efficient time synchronous as compare to FTSP (Flooding Time Synchronization protocol). In traditional protocol common strategy is to flood current time information of a reference node into the network. Alternate method of same is proposed which allow node to propagate as soon as possible but this scheme is difficult have many limitation such as energy and network performance, in this paper author proposed a method to reduce the undesired effect of on synchronization accuracy without altering the propagation speed of the flood, main objective is to enhance the accuracy and stability of protocol SFTP drastically.

III. PROPOSED METHOD

In this module we review the various requirement for synchronization and tradeoff between the efficiency and the accuracy (e.g., precision and energy efficiency), can say that scheme cannot balance them in combination. A synchronization method must measures accurately with increase in the number of sensor node. The need for accuracy and efficiency may vary or dependent on the type of application and the reason for synchronization. To achieve this task sensor need to agree on communication criteria such as time, traditional method of clock synchronization is not appropriate may be due to huge complexity and more energy requirement. For instance NTP protocol work adequately on internet to synchronize systems, but same protocol not suitable for WSN's also GSP could be expansive to fix in low costing devices. Network time synchronization methods are used in Computer network techniques described in late 1981 (IEN) 173 and a public protocol was developed from it that was documented in RFC 778[16][4], complete specification of the NTPv1 protocol algorithms, was exhibit in RFC 1059. A typical NTP client will regularly poll three or more servers on diverse networks. To synchronize clock with a server, the client required to calculate the RTT delay and the offset[4][16]. Consider the client sever mode and formula to calculate the RTT is given in equation 2, t_0 is the timestamp of the client request packet transmission, t_1 is the server's timestamp of the request packet reception, t_2 is the server's timestamp of the response packet transmission and t_3 is the client's timestamp of the response packet reception also we can calculate the offset denoted by θ in equation 3, to compute the θ , calculate $t_3 - t_0$ is the time elapsed on the client side between the transmission of the request packet and the reception of the response packet and $t_2 - t_1$ is the time the server waited before sending the reply, the offset θ is given by equation 3[13], The synchronization problem consists of four parts as shown in figure 3.1, as exhibited send time, access time, propagation time, and receive time. Calculated δ and θ are passed through filters and subjected to statistical analysis Method to improve the robustness and scalability; many distributed protocols are designed for TSP in Wireless Sensor Network. Solis, et al. [6] proposed a complete asynchronous distributed time synchronization protocol (DTSC). Giridhar proposed DTSC as a coordinate-descent optimization problem[17]. Currently, the agreement based method has been designed for TS in Wireless sensor network. A agreement method has three major advantages over conventional method, able to work in spatially distributed way, so would not require a tree topology or a root node as a server to synchronized [6]. Second, depends on the agreement algorithm, more precision clocks between nodes, and lastly it helps to remunerate the skew and offset differences between nodes, where as to calculate d propagation delay of time as below.

$$\theta = ((t_1 - t_0) - (t_2 - t_3)) / 2$$

$$d=((t_1-t_0) + (t_3-t_2))/2$$

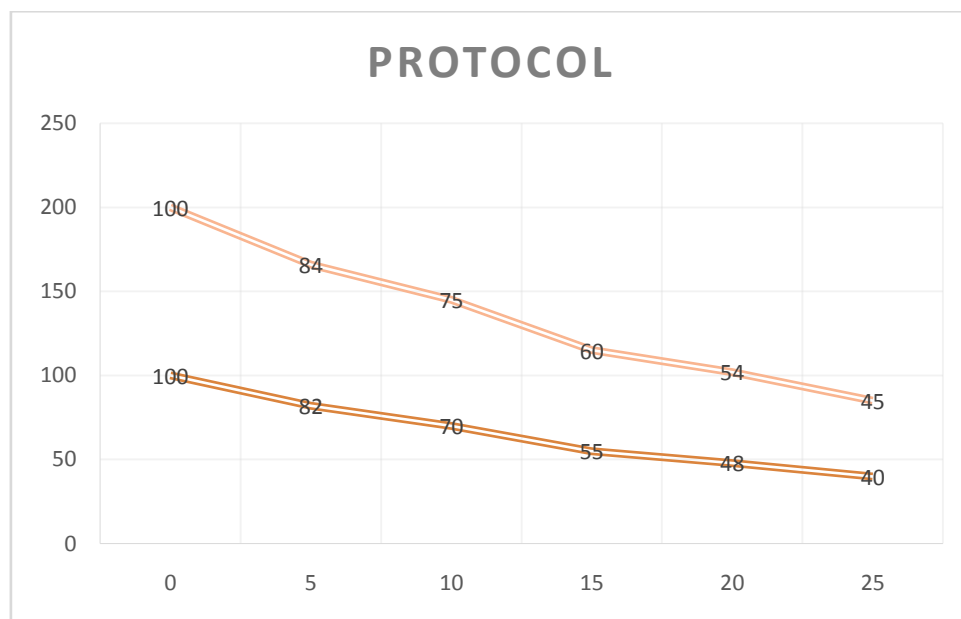
4

We propose that synchronization overhead should be reduces wherever possible randomly. The number of messages required in synchronization especially in the case of always-active type of network, is much more than even the messages sent for notification of events. This method is not suitable in all the scenarios but it provides the necessary and sufficient synchronization to the scenario where numbers of events are very low.

For selecting the root node we make use of election algorithm for periodic root node election step. After selection of root node may assigned as level 0, and generate the hierarchy level for rest of the nodes, discovery phase by broadcasting a level discovery packet [14]. This packet comprise the identity and level of the sender node. Upon receiving this packet, the neighbors of the root node allot themselves level 1. Each node in level 1 will ruckus the level discovery packet including level and identity in the packet. Later fixing the level of each node no further level discovery packet is entertained, which is already achieved by chain broadcasting method between the nodes

In the proposed work the simulation of the protocols is done using a network of wireless sensor network. This network consists of a sink node and various numbers of sensors. Cluster heads are also simulated as per their use in our algorithm. As discussed in the previous chapter, Cluster head are chosen using the already existing protocols. For the purpose of choosing cluster head, any of the existing method like LEACH can be used[9]. Choosing the cluster head put a lot of overhead in our network, but we do not consider this overhead in our simulation results. Non deterministic delays are calculated using distributions and range of delays are taken from Table 1. The message that is used for synchronization is of same length in this simulation work. The delays that are discussed and used in our simulation are taken from the article, in which the author study and analysis the results using hardware sensor node i.e mica motes.

For selecting the root node we make use of election algorithm for periodic root node election step. After selection of root node may assigned as level 0, and generate the hierarchy level for rest of the nodes, discovery phase by broadcasting a level discovery packet [14]. This packet comprise the identity and level of the sender node. Upon receiving this packet, the neighbors of the root node allot themselves level 1. Each node in level 1 will ruckus the level discovery packet including level and identity in the packet. Later fixing the level of each node no further level discovery packet is entertained, which is already achieved by chain broadcasting method between the nodes



Conclusion

A conclusion section must be included and should indicate clearly the advantages, limitations, and possible applications of the paper. Although a conclusion may review the main points of the paper, do not replicate the abstract as the conclusion. A conclusion might elaborate on the importance of the work or suggest applications and extensions. Significance of time synchronization protocol in wireless sensor network is expounded and compared with other traditional Time Synchronization Protocols. Extensive literature survey was carried out and found measure of performance in single protocol was not competent for different wireless network subsequently required variant approach. Divergence in measure of performance may lead to stipulate in services accordingly. We acquaint the rising trend in the WSN's and design a method to mitigate the anarchy in

the network due to dissemination of messages in the network during flooding. Flooding is one of the major protocol at the beginning of the routing technique in wireless sensor network, which validates less complexity but peak in consumption of resources. Various authors use these protocol for the communication between the source and destination. Many review paper predict that simply forwarding the packets after regular interval using flooding technique can lead to poor performance in context to energy, bandwidth, and precision. It was also shown that if we decrease the propagation speed of the slow flooding lead to decrease in synchronization accuracy and stability. Alternate solution to this problem is rapid flooding indeed incorporate the accuracy but may decrease in the network life time, our task in this network is explicitly associated with to maintain the stability in the network with minimum packet generation. Energy required in proposed protocol as compared to traditional protocol is effectively less. Our theoretical finding and experimental results exhibits that the proposed scheme mitigate the unnecessary forwarding of packets to neighbor node, also more the neighbor near the sender more number of packet is need to generate. Employing agnostic approach and considering the good aspects of the RBS and Slow flooding drastically improve the precision and stability in the network. In this context we postulates that the hybrid technique proposed is pragmatic and sufficient to ameliorate the existing latency error in synchronization of clock and slow drift in clock. The role of a sensor can be classified as sensing and routing. Sensor predict the path by collecting the information from nearby sensor, depending upon the hierarchy and protocol used in the network. All the operations are battery consuming and saving battery is one of the main concerned of the researchers. The algorithm provide minimum overhead of message transmission and hence resulted into ample energy save at each sensor node. This will also increase the network life time or prolong network. Simulated the proposed Hybrid method and comparing it with the individual protocols on same parameter shows that the Hybrid method provides low overhead as compared to other methods.

REFERENCES

-
- [1]. J. Wu, L. Zhang, Y. Bai, and Y. Sun, "Cluster-Based Consensus Time Synchronization for Wireless Sensor Networks," vol. 15, no. 3, pp. 1404–1413, 2015.
 - [2]. V. Dhiman, V. Kaur, and A. Singh, "Performance investigation of energy aware routing protocol for heterogeneous WSNs."
 - [3]. F. Sivrikaya and B. Yener, "Time synchronization in sensor networks: a survey," *IEEE Netw.*, vol. 18, no. 4, pp. 45–50, 2004.
 - [4]. J. Elson, L. Girod, and D. Estrin, "Fine-grained network time synchronization using reference broadcasts," *ACM SIGOPS Oper. Syst. Rev.*, vol. 36, no. SI, p. 147, 2002.
 - [5]. K. S. Yildirim and A. Kantarci, "Time synchronization based on slow-flooding in wireless sensor networks," *IEEE Trans. Parallel Distrib. Syst.*, vol. 25, no. 1, pp. 244–253, 2014.
 - [6]. S. Hwang and Y. Baek, "Reliable time synchronization protocol for wireless sensor networks," *Embed. Ubiquitous Comput. - Euc 2005*, vol. 3824, no. August, pp. 663–672, 2005.
 - [7]. M. Windows, M. Os, C. P. When, Y. Wei, P. Yildirim, C. den Bulte, C. Dellarocas, T. Weekly, I. C. T. I. Weekly, W. E. Henley, S. D. Vyas, T. Uk, S. Trend, F. S. Trend, B. Technology, F. Insights, A. Longtop, F. Technologies, A. Tan, S. Darby, I. O. I. Corp, K. Loong, SW공학센터, J. Skan, R. Lumb, S. Masood, S. K. Conway, K. Shue, S. N. Service, W. Server, S. Module, F. Security, L. Scheme, C. Sample, K. Schaffer, S. Report, I. I. Report, M. Pos, A. Permissions, W. Experience, A. Links, M. Payments, F. Support, P. Store, W. Paper, F. I. Online, C. Okten, U. O. Osili, S. F. November, L. Name, F. Name, O. Training, P. Training, C. Darin, R. O. Training, M. Kimberly, G. Deepa, E. Board, E. Principal, I. Primary, F. Systems, E. B. Study, N. Co-investigator, R. Mohamad, A. Building, N. A. Ismail, S. F. March, M. Lin, N. R. Prabhala, S. Viswanathan, K. Lee, S. Park, J. J. Lee, S. Park, F. Law, R. J. Straight, S. Vice, C. Privacy, U. Straight, D. B. Douglass, B. Y. C. Avery, G. Fanger, D. B. Douglass, KPMG, D. Kempe, J. Kleinberg, É. Tardos, C. Karma, S. Issues, M. Issue, S. Internet, P. Service, L. G. B. Insight, H. Indicators, C.-L. Huang, M.-C. Chen, C.-J. Wang, A. Group, R. P. F. Go, R. P. F. Go, D. D. Go, D. D. Go, R. P. F. Go, I. O. From, S. Freedman, G. Jin, E. P. Forgot, B. D. Foust, A. P. February, F. I. Corp, B. Fair, E. Isaac, M. Fellowes, F. Isaac, T. Fico, F. Isaac, A. B. Sanders, D. Bank, N. York, W. Street, T. Fico, E. For, T. H. E. Home, S. Finance, S. F. February, M. Eroglu S., Toprak S., Urgan O, MD, Ozge E. Onur, MD, Arzu Denizbasi, MD, Haldun Akoglu, MD, Cigdem Ozpolat, MD, Ebru Akoglu, Ernst & Young, S. S. M. Economics, J. Duarte, S. Siegel, L. Young, S. Debnath, N. Ganguly, P. Mitra, B. Data, T. S. Dapp, M. L. T. Cossio, L. F. Giesen, G. Araya, M. L. S. Pérez-Cotapos, R. L. VERGARA, M. Manca, R. A. Tohme, S. D. Holmberg, T. Bressmann, D. R. Lirio, J. S. Román, R. G. Solís, S. Thakur, S. N. Rao, E. L. Modelado, A. D. E. La, C. Durante, U. N. A. Tradición, M. En, E. L. Espejo, D. E. L. A. S. Fuentes, U. A. De Yucatán, C. M. Lenin, L. F. Cian, M. J. Douglas, L. Plata, F. Héritier, T. G. Contributor, M. S. Conference, S. M. Call, G. Compass, B. Cho, J. Park, J. Chavan, W. Chat, O. Channel, B. Challenge, K. Brief, K. Brief, B. Block, S. S. M. Economics, U. Big, D. Based, N. Analysis, A. Bankitx, C. Bank, C. Bank, C. Bank, B. Bankers, H. Street, B. B. Bank, F. C. Authority, and Accenture, *No Title No Title*, vol. XXXIII, no. 2. 2014.
-

- [8]. M. Maruti, B. Kusy, G. Simon, and L. Kos L. deczi, "The flooding time synchronization protocol," *Proc. 2nd Int. Conf. Embed. networked Sens. Syst. - SenSys '04*, p. 39, 2004.
- [9]. B. K. Debroy, M. S. Sadi, and A. Imran, "An Efficient Approach to Select Cluster Head in Wireless Sensor Networks," vol. 6, no. 7, pp. 529–539, 2011.
- [10]. Virpalkaur and Vikramdhiman, "A SURVEY OF ENERGY CONSUMPTION ROUTING PROTOCOLS IN WIRELESS SENSOR NETWORKS," vol. 15, 2015.
- [11]. J. Hill and D. Culler, "A wireless embedded sensor architecture for system-level optimization."
- [12]. J. Elson and D. Estrin, "Time Synchronization for Wireless Sensor Networks," *Proc. 2001 Int. Parallel Distrib. Process. Symp. (IPDPS), Work. Parallel Distrib. Comput. Issues Wirel. Networks Mob. Comput.*, pp. 1965–1970, 2001.
- [13]. J. Elson and D. Estrin, "Time Synchronization for Wireless Sensor Networks."
- [14]. B. Sundararaman, U. Buy, and A. D. Kshemkalyani, "Clock Synchronization for Wireless Sensor Networks: Survey," *Ad Hoc Networks*, vol. 3, no. 3, pp. 281–323, 2005.
- [15]. D. L. Mills and F. Group, *Network The Network Time Protocol*. 2006.
- [16]. D. L. Mills, F. Acm, and F. Ieee, "A Brief History of NTP Time : Memoirs of an Internet Timekeeper 1," pp. 1–13, 1981.
- [17]. J. He, P. Cheng, L. Shi, J. Chen, and Y. Sun, "Time synchronization in WSNs: A maximum-value-based consensus approach," *IEEE Trans. Automat. Contr.*, vol. 59, no. 3, pp. 660–675, 2014.
- [18]. A. Nayyer, M. Nayyer, and L. K. Awasthi, "A Comparative study of Time Synchronization Protocols in Wireless Sensor Network," *Int. J. Comput. ...*, vol. 36, no. 11, pp. 13–19, 2011.
- [19]. A. Varga, "THE OMNET ++ DISCRETE EVENT SIMULATION SYSTEM," 2000.
- [20]. E. Mahalakshmi and A. Rajamurugan, "Reference Broadcast Time Synchronization in WSN," vol. 3, no. 6, pp. 17–20, 2015.